

Introduction to MPI Programming

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Outline

- What is the MPI library?
 - How to compile it in your program;
 - Basic commands and interfaces;
 - Sample programs;

What will you learn:

Basic information necessary to using an MPI library, and start self-learning from there.

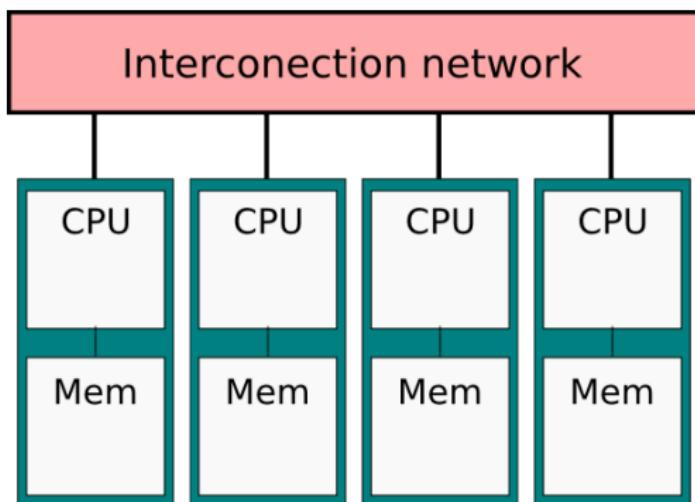
Download the materials from this class (notes and code)

<http://conclave.cs.tsukuba.ac.jp/pub/MPIintro.zip>



Distributed Memory Machine

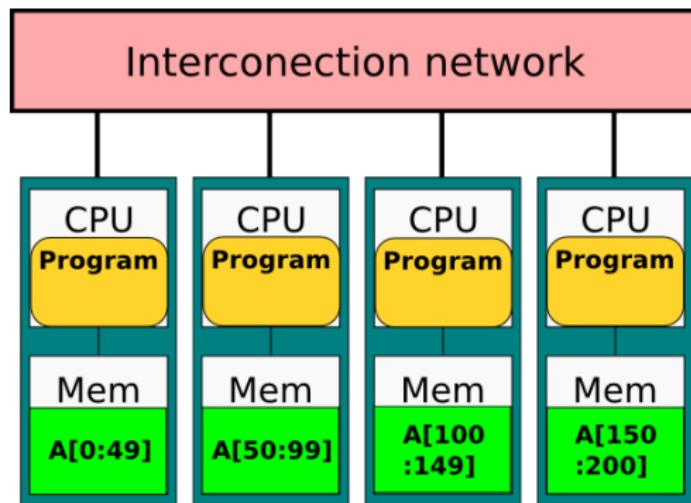
- Computers/Nodes, connected by a interconnection network;
 - Node: a CPU and memory;
 - Each node execute a program, and communicates data through the network.





Single Program, Multiple Data (SPMD)

- Parallel of the same program independently;
 - Data is different for each instance;
 - Programs interact through message exchange;





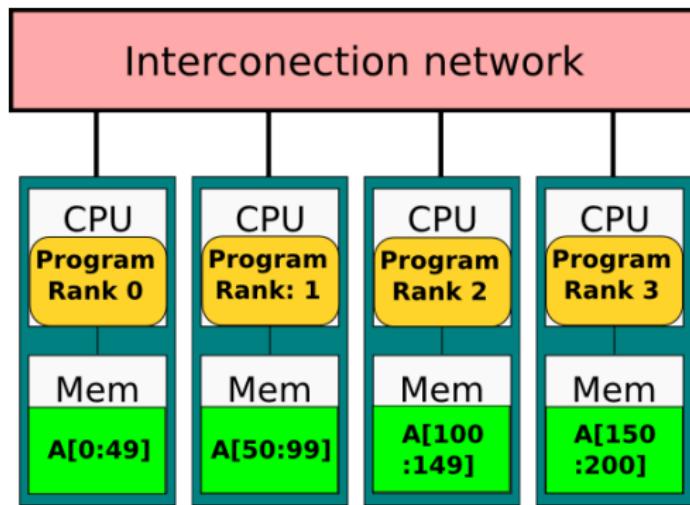
MPI - The Message Passing Interface

- Standard for Message Passing Interface;
 - Portable Parallel Library;
 - 8 communication modes, collective communication, communication domain, process topology;
 - Defines more than 100 interfaces;
 - C, C++, Fortran (and many wrappers);
 - Frequent Updates by the MPI Forum:
(<http://www.mpi-forum.org>)
 - MPI-1 (1994)
 - MPI-2 (2009)
 - MPI-3 (2014)
 - MPI-4 (in Discussion)

MPI execution model



- Execute the same program on each processor;
- Not synchronous if no communication happens;
- Each process has an ID (rank);





Requirements of an MPI program: code requirements

These are the essential headers for an MPI program in C:

```
#include <mpi.h> /* MPI library headers */  
int main(int argc, char *argv[])  
{  
    MPI_Init(&argc,&argv)  
    /* Must be called before ANY MPI functions */  
  
    MPI_Finalize()  
    /* Must be called before the end of the program */  
}
```

Note that all MPI functions and structure are prefixed by MPI_



Compiling an MPI Program

- **specialized compiler:** `mpicc -Ilibs source`
MPI library itself does not need to be included.
- **specialized run:** `mpiexec -n N binary`
N is the number of processes
For other options, see the man page

MPI libraries

Some packages with MPI implementations on ubuntu 14.04:

- OpenMPI: `openmpi-bin, libopenmpi-dev`
- MPICH: `mpich, libmpich-dev`

Program: sample1.c



```
#include <stdio.h>
#include <mpi.h>

int main(int argc, char *argv[])
{
    int rank,len;
    char name[MPI_MAX_PROCESSOR_NAME] ;

    MPI_Init(&argc,&argv);
    MPI_Comm_rank(MPI_COMM_WORLD,&rank);
    MPI_Get_processor_name(name,&len);
    printf("Processor %03d: %s; reporting!\n", rank, name);

    MPI_Finalize();
    return(0);
}
```

Program: sample1.c



- Each processor runs the same program, but has different rank numbers, and possibly processor names (local memory);
- Rank numbers are defined in relation to the Communicators (COMM_WORLD);
- Use Init and finalize functions around MPI functions;



Communicator (2)

Communication Domain

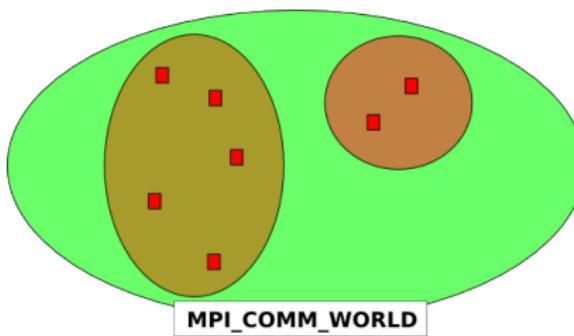
- Group of Processes;
- May be defined around a topography (1D Ring, 2D Mesh, Torus, etc);
- Stores number of processes, Rank for each process;

`MPI_COMM_WORLD`: Initial communicator, including all processes;



Communicator (2)

- Communicators define “scopes” for collective communication; e.g.: two thirds of the processes calculate the weather forecast for the current weather forecast; one third compute the initial conditions for the next iteration;
- Intra-communicator and Inter-communicator communication;





Communicator (3)

```
int MPI_Comm_size(MPI_Comm comm, int *size)
```

Returns the total number of processes in the communicator `comm`;

```
int MPI_Comm_rank(MPI_Comm comm, int *rank)
```

Returns the process rank in the communicator `comm`;



Collective Communication

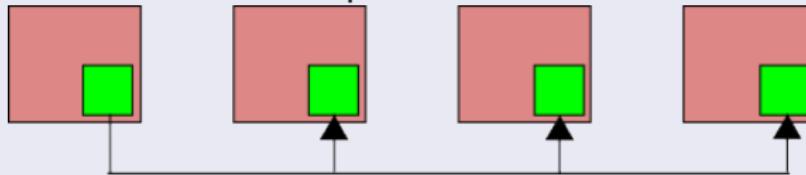
- Refers to communication that involves all processes (of a communicator);
- Barrier Synchronization (depends on implementation);
- Global Communication: Broadcast, Gather, Scatter, allgather, allscatter, etc;
- Global Reduction: Reduction (sum, maximum, local operator), Prefix computations;



Simple Collective Communication

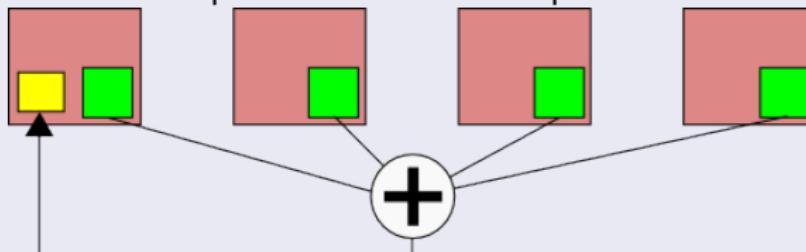
Broadcast

Sends data from one process to all.



Reduce

Gathers and process data from all processes.



- Library optimizes the implementation based on the architecture (no work for the programmer).



MPI_Bcast

```
int MPI_Bcast(void *buffer, int count, MPI_Datatype type,  
              int root, MPI_Comm comm);
```

- returns error value;
- **buffer**: starting address of data;
- **count**: number of units in the data;
- **type**: data type of a data unit;
- **root**: which rank is sending the data;
- **comm**: communicator;



MPI_Reduce

```
int MPI_Reduce(void *sendbuf, void* recvbuf, int count,  
               MPI_Datatype type, MPI_Op operator,  
               int root, MPI_Comm comm);
```

- **sendbuf**: data to be sent
- **recvbuf**: resulting data
- **operator**: reduce operation



Program: sample2.c

Parallel Summation

- Serial Summation:

$$A_1 + A_2 + A_3 + \dots + A_{999} = \text{sum}$$

- Parallel Summation:

$$A_1 + A_2 + A_3 + \dots + A_{333} = B_1$$

$$A_{334} + A_{335} + A_{336} + \dots + A_{666} = B_2$$

$$A_{667} + A_{668} + A_{669} + \dots + A_{999} = B_3$$

$$B_1 + B_2 + B_3 = \text{sum}$$

- Each process has a different part of the data
- Calculate its own sum, and put everything together with *Reduce*



Program: sample2.c

```
/**  
 * Sample program 2 - sum a large array of numbers  
 */  
  
#include <stdio.h>  
#include <mpi.h>  
#include <string.h>  
  
int main(int argc, char *argv[])  
{  
  
    int local_tsum = 0;  
    int local_psum = 0;  
    int local_rank = 0;  
    int local_data[250];  
  
    int local_i;
```



Program: sample2.c

```
MPI_Init(&argc,&argv);
MPI_Comm_rank(MPI_COMM_WORLD,&local_rank);

/* Initializing some toy data */
for(local_i = 0; local_i < 250; local_i++)
    local_data[local_i] = local_rank;
/* You would have the real data here */

for(local_i = 0; local_i < 250; local_i++)
    local_psum += local_data[local_i];

MPI_Reduce(&local_psum,&local_tsum,1,MPI_INT,MPI_SUM,
           0,MPI_COMM_WORLD);
```



Program: Sample2.c

```
printf("Sum for process rank %d: %d\n", local_rank,  
      local_psum);  
fflush(0);  
  
MPI_Barrier(MPI_COMM_WORLD);  
  
if (local_rank == 0)  
    printf("Total Sum: %d\n", local_tsum);  
  
MPI_Finalize();  
return(0);  
}
```



Program: sample3.c

Pi Calculation

- Use the Riemann Sum;
$$\sum_0^1 \frac{4}{1+t^2} \Delta t$$
- Choose n (number of divide parts) and broadcast;
- Each processor calculates a number of parts;
p1: 1,4,7,10...n;
p2: 2,5,8,11...n+1;
p3: 3,6,9,12...n+2;
- Partial sums are reduced;

Hint: Try to execute with $N > 64000$ to see the speed difference!



Program: sample3.c

```
/**  
 * Sample program 3 - calculate PI using the Riemann Sum  
 */  
  
#include <stdio.h>  
#include <stdlib.h>  
#include <mpi.h>  
#include <math.h>  
  
double f(double a)  
{  
    return (4.0/(1.0 + a*a));  
}
```



Program: sample3.c

```
int main(int argc, char *argv[])
{
    int n = 0, myid, nprocs, i;
    double PI25DT = 3.141592653589793238462643;
    double mypi, pi, h, sum, x;
    double startwtime = 0.0, endwtime;

    int namelen;
    char processor_name[MPI_MAX_PROCESSOR_NAME];

    MPI_Init(&argc, &argv);
    MPI_Comm_size(MPI_COMM_WORLD, &nprocs);
    MPI_Comm_rank(MPI_COMM_WORLD, &myid);
    MPI_Get_processor_name(processor_name, &namelen);

    if (argc > 1)
        n = atoi(argv[1]);
    startwtime = MPI_Wtime();
```



Program: sample3.c

```
/* broadcast 'n' */
MPI_Bcast(&n, 1, MPI_INT, 0, MPI_COMM_WORLD);
if (n <= 0)
{
    fprintf(stderr, "usage: %s <#partition>\n", argv[0]);
    MPI_Abort(MPI_COMM_WORLD,1);
}

/* calculate each part of pi */
h = 1.0/n;
sum = 0.0;
for (i = myid+1; i <= n; i+= nprocs)
{
    x = h * (i-0.5);
    sum += f(x);
}
mypi = h * sum;
```



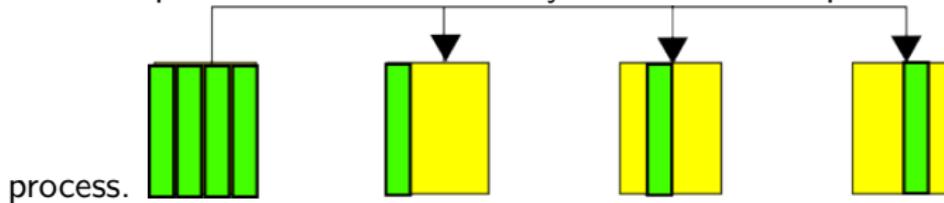
Program: sample3.c

```
/* sum up each part of pi */  
MPI_Reduce(&mypi, &pi, 1, MPI_DOUBLE, MPI_SUM, 0,  
            MPI_COMM_WORLD);  
  
if (myid == 0)  
{  
    printf("pi is approximately %.16f, Error is %.16f\n",  
          pi, fabs(pi - PI25DT));  
    endwtime = MPI_Wtime();  
    printf("wall clock time = %f\n",  
          endwtime - startwtime);  
}  
MPI_Finalize();  
return(0);  
}
```



A few more collective communication functions

- **scatter:** split the data into subarrays and send each part to a



- **gather:** get split data from each process and put it together in an

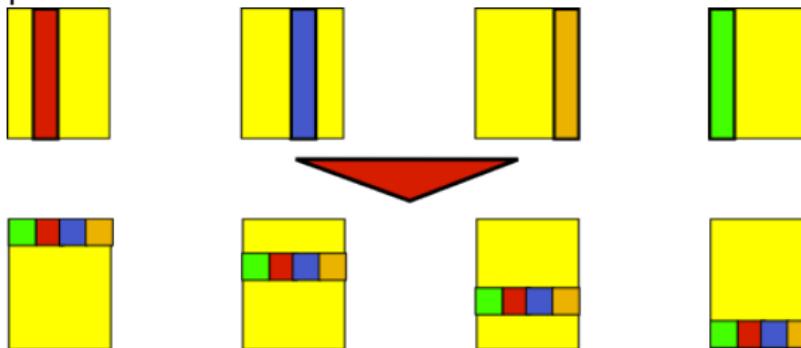




A few more collective communication functions

"All-*" functions.

- **allreduce**: result of reduce operation is sent to all processes
- **allgather**: split data is put together and sent to all processes
- **allscatter, alltoall**: data at each process is split and sent to other processes



(can be seen as a matrix transformation of 2D data);



Point to Point communication

Data Transfer Between Two Processes

- Process A sends some data to process B (send);
- Process B receives the data from process A (recv);



Point to Point communication (2)

- Data Type specified in MPI Call;
Basic type, array, structure, vector, user-defined;
(MPI_INT, MPI_DOUBLE,...);
- The send/receive pair is specified by:
Communicator, **message tag**, source rank, destination rank;
(can use some wildcards: MPI_ANY_TAG, MPI_ANY_SOURCE, etc);



Blocking and Non-Blocking Communication

Blocking

- `Send(A...)` returns when send buffer can be re-used;
But the message isn't transmitted yet.
- `Receive` returns when the receive buffer is available for use;

Non-Blocking

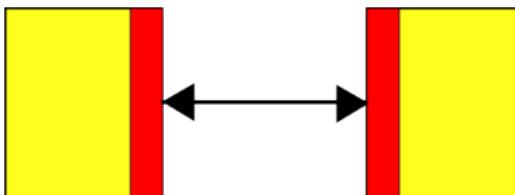
- `Send` and `Receive` are separated into “post” and “complete”;
- Enables overlapping of computation and communication;

Note that it is perfectly legal to mix both!

Blocking and Non-Blocking Communication (2)



Message Exchange



Blocking

- `MPI_Send(dest,data)`
- `MPI_Recv(src,data)`

If both sides do this at the same time, in a **non-buffered** mode, may result in a **deadlock**

Non-Blocking

- `MPI_Isend(dst,data,request)`
- `MPI_Irecv(src,data,request)`
- ...
- `MPI_Waitall(request)`

Message exchange always complete, regardless of comm mode. Portable.



Communication Modes

4 Communication modes dictate the behavior of send and receive:

- **Standard mode:** MPI decides if the message should be buffered or not. Can't assume it is buffered!
- **Buffered mode:** Outgoing message is always buffered. Send operation guaranteed to be local;
User must guarantee the buffer!
- **Synchronous mode:** Send completes only if a matching receive is posted, and data begins transferring.
Send operation guaranteed to be non-local;
- **Ready mode:** Checks if a receive has been posted before starting send. Sending without a matching receive return error.
Can remove hand-shake operations;

Large combination of communication modes, blocking and non blocking sends and receives. **Check the manual!**.



Blocking Communication

MPI_Send

```
MPI_Send(void *data, int count, MPI_Datatype type,  
        int dest, int tag, MPI_Comm comm);
```

MPI_Recv

```
MPI_Recv(void *data, int count, MPI_Datatype type,  
         int dest, int tag, MPI_Comm comm,  
         MPI_Status *status);
```

- **Tag** - message identifier
- **Status** - structure with tag, source and other data.



Non Blocking Communication

MPI_Isend

```
MPI_Isend(void *data, int count, MPI_Datatype type,  
          int dest, int tag, MPI_Comm comm,  
          MPI_Request *request);
```

MPI_Irecv

```
MPI_Irecv(void *data, int count, MPI_Datatype type,  
          int dest, int tag, MPI_Comm comm,  
          MPI_Request *request);
```

MPI_Wait

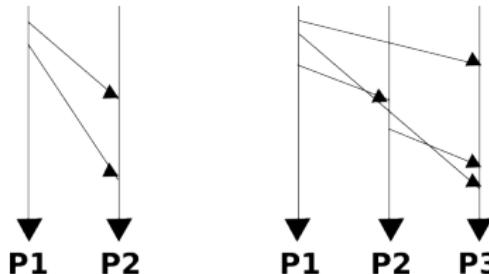
```
MPI_Wait(MPI_Request *request, MPI_Status *status);
```

- Request: Contains information about the message received or sent.

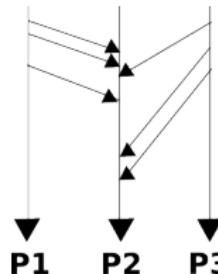


Communication Caveats

- **Caveat 1:** With more than 2 processes, message arrival order is not guaranteed;



- **Caveat 2:** Message fairness is not guaranteed either;

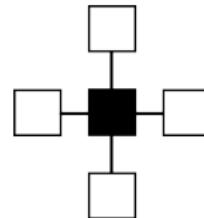




Program: sample4.c

Explicit Solution of Laplace equation

- update by averaging data of four neighbor cells;
- two arrays with old and new values;
- segmentation by region;
- compute residual to check for convergence;

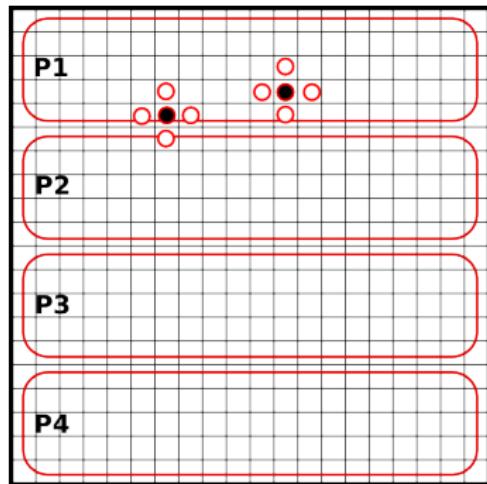
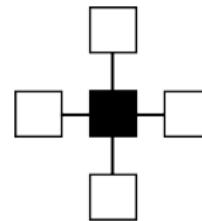


Program: sample4.c



Explicit Solution of Laplace equation

- Block distribution of the 2D data;
- Boundary elements require neighborhood communication;
- Data exchange between boundary elements (Arrays);





Process Topology

MPI can create a **virtual** spatial relationship between communicators (above, below, etc), called a topology.

```
int MPI_Cart_create(MPI_Comm comm_old, int ndims,  
                    int *dims, int *periods, int reorder,  
                    MPI_Comm *comm_cart);
```

- **ndims**: number of dimensions in the grid;
- ***dims**: size for each dimension;
- ***periods**: whether each dimension wraps around;
- **reorder**: if the ranks can be reordered for optimization;



Process Topology

To calculate the *dest*, we use a `Cart_shift` function.

```
int MPI_Cart_shift(MPI_Comm comm, int direction,  
                   int disp, int *rank_source,  
                   int *rank_dest);
```

- `comm`: Holds the grid topology;
- `direction`: dimension of the shift (horizontal, vertical, etc);
- `disp`: size of the shift;
- Returns `MPI_PROC_NULL` rank if boundary is non-periodical, and tries to go to an invalid rank;



Program: sample4.c

```
/**  
 * Sample program 4 - laplace equation with explicit method  
 */  
  
#include <stdio.h>  
#include <stdlib.h>  
#include <math.h>  
#include <mpi.h>  
  
/* Square Region */  
  
#define XSIZE 256  
#define YSIZE 256  
#define PI 3.1415927  
#define NITER 10000  
  
double u[XSIZE+2] [YSIZE+2], uu[XSIZE+2] [YSIZE+2];  
/* 2d target domain, uu is used for new values */
```





Program: sample4.c

```
double time1, time2;
void lap_solve(MPI_Comm);
int myid, numprocs;
int namelen;
char processor_name[MPI_MAX_PROCESSOR_NAME];
int xszie;

void initialize()
{
    int x,y;

    /* initialization */
    for (x=1; x < XSIZE + 1; x++)
        for (y=1; y < YSIZE + 1; y++)
            u[x][y] = sin((x - 1.0)/XSIZE*PI) +
cos((y-1.0)/YSIZE*PI);
```



Program: sample4.c

```
/* zero clear in the boundary */
for (x = 0; x < XSIZE + 2; x++)
{
    u[x][0] = u[x][YSIZE+1] = 0.0;
    uu[x][0] = uu[x][YSIZE+1] = 0.0;
}
for (y = 0; y < YSIZE + 2; y++)
{
    u[0][y] = u[XSIZE+1][y] = 0.0;
    uu[0][y] = uu[XSIZE+1][y] = 0.0;
}
#define TAG_1 100
#define TAG_2 101
#ifndef FALSE
#define FALSE 0
#endif
```



Program: sample4.c

```
void lap_solve(MPI_Comm comm)
{
    int x,y,k;
    double sum, t_sum;
    int x_start, x_end;
    MPI_Request req1, req2;
    MPI_Status status1, status2;
    MPI_Comm comm1d;
    int down, up;
    int periods[1] = { FALSE };

    /* Create one dimensional cartesian topology with non
     * periodical boundary */
    MPI_Cart_create(comm, 1, &numprocs, periods,
                   FALSE, &comm1d);
    /* calculate process ranks for 'down' and 'up' */
    MPI_Cart_shift(comm1d,0,1,&down,&up);
```



Program: sample4.c

```
x_start = 1 + xszie*myid;
x_end = 1 + xszie*(myid+1);

for (k = 0; k < NITER; k++)
{ /* old <- new */
    for (x = x_start; x < x_end; x++)
        for (y = 1; y < YSIZE+1; y++)
            uu[x][y] = u[x][y];

    /* recv from down */
    MPI_Irecv(&uu[x_start-1][1], YSIZE, MPI_DOUBLE, down,
              TAG_1, comm1d, &req1);
    /* recv from up */
    MPI_Irecv(&uu[x_end][1], YSIZE, MPI_DOUBLE, up, TAG_2,
              comm1d, &req2);
```



Program: sample4.c

```
/* send to down */
MPI_Send(&u[x_start] [1] ,YSIZE,MPI_DOUBLE,down,TAG_2,
comm1d);
/* send to up */
MPI_Send(&u[x_end -1] [1] ,YSIZE,MPI_DOUBLE,up,TAG_1,
comm1d);

MPI_Wait(&req1,&status1);
MPI_Wait(&req2,&status2);

/* update */
for (x = x_start; x < x_end; x++)
    for (y = 1; y < YSIZE + 1; y++)
        u[x] [y] = .25*(uu[x-1] [y] + uu[x+1] [y] +
                         uu[x] [y-1] + uu[x] [y+1]);
}
```

Program: sample4.c



```
/* check sum */  
sum = 0.0;  
for (x = x_start; x < x_end; x++)  
    for (y = 1; y < YSIZE + 1; y++)  
        sum += uu[x][y] - u[x][y];  
MPI_Reduce(&sum, &t_sum, 1, MPI_DOUBLE, MPI_SUM, 0,  
           comm1d);  
if (myid == 0) printf("summ = %g\n", t_sum);  
MPI_Comm_free(&comm1d);  
}  
int main(int argc, char *argv[])  
{  
    MPI_Init(&argc, &argv);  
    MPI_Comm_size(MPI_COMM_WORLD, &numprocs);  
    MPI_Comm_rank(MPI_COMM_WORLD, &myid);  
    MPI_Get_processor_name(processor_name, &nnameLEN);  
    fprintf(stderr, "Process %d on %s reporting for duty.\n",  
           myid, processor_name);
```



Program: sample4.c



```
xsize = XSIZE/numprocs;
if ((XSIZE % numprocs)!=0)
    MPI_Abort(MPI_COMM_WORLD,1);

initialize();
MPI_Barrier(MPI_COMM_WORLD);

time1 = MPI_Wtime();
lap_solve(MPI_COMM_WORLD);
MPI_Barrier(MPI_COMM_WORLD);
time2 = MPI_Wtime();

if (myid == 0)
    printf("time = %g\n", time2 - time1);
MPI_Finalize();

return(0);
}
```





Things to Improve

- This program allocates the whole array, although it is not necessary
 - When a partial array is allocated, take care about computing the local and the global indexes.
 - This is essential for large scale problems, using a distributed memory machine.

- Two dimensional distribution of 2D array is more efficient than one dimensional distribution
 - Reduces the communication size
 - Can be parallelized with a larger number of processors



Open Source MPI

- OpenMPI:
<http://www.open-mpi.org/>
- MPICH2:
<http://www-unix.mcs.anl.gov/mpi/mpich2/>
- YAMPII:
<http://ww.il.is.s.u-tokyo.ac.jp/yampii/>