



Accelerating Fermionic Molecular Dynamics

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Abstract

We consider how to accelerate Fermionic Molecular Dynamics algorithms by introducing n pseudofermion fields coupled with the n^{th} root of the Fermionic kernel. The n^{th} roots may be computed efficiently using Чебышев rational approximations, and to this end we review the theory of optimal polynomial and rational Чебышев approximations.



Non-linearity of CG solver

- ▶ Suppose we want to solve $A^2x=b$ for Hermitian A by CG
 - ▶ It is better to solve $Ax=y$, $Ay=b$ successively
 - ▶ Condition number $\kappa(A^2) = \kappa(A)^2$
 - ▶ Cost is thus $2\kappa(A) < \kappa(A^2)$ in general
- ▶ Suppose we want to solve $Ax=b$
 - ▶ Why don't we solve $A^{1/2}x=y$, $A^{1/2}y=b$ successively?
- ▶ The square root of A is uniquely defined if $A > 0$
 - ▶ This is the case for fermion kernels
- ▶ All this generalises trivially to n^{th} roots
 - ▶ No tuning needed to split condition number evenly
- ▶ How do we apply the square root of a matrix?



Rational matrix approximation

- ▶ Functions on matrices
 - ▶ Defined for a Hermitian matrix by diagonalisation
 - ▶ $H = UDU^{-1}$
 - ▶ $f(H) = f(UDU^{-1}) = U f(D) U^{-1}$
- ▶ Rational functions do not require diagonalisation
 - ▶ $\alpha H^m + \beta H^n = U(\alpha D^m + \beta D^n) U^{-1}$
 - ▶ $H^{-1} = UD^{-1}U^{-1}$
- ▶ Rational functions have nice properties
 - ▶ Cheap (relatively)
 - ▶ Accurate



Polynomial approximation

- ▶ What is the best polynomial approximation $p(x)$ to a continuous function $f(x)$ for x in $[0, 1]$?
 - ▶ Best with respect to the appropriate norm

$$\|p - f\|_n = \left(\int_0^1 dx |p(x) - f(x)|^n \right)^{1/n}$$

where $n \geq 1$

Weierstraß' theorem

- ▶ Taking $n \rightarrow \infty$ this is the “minimax” norm

$$\|p - f\|_{\infty} = \min_p \max_{0 \leq x \leq 1} |p(x) - f(x)|$$

- ▶ Weierstraß: Any continuous function can be arbitrarily well approximated by a polynomial



Бернштейне polynomials



► The explicit solution is provided by Бернштейне polynomials

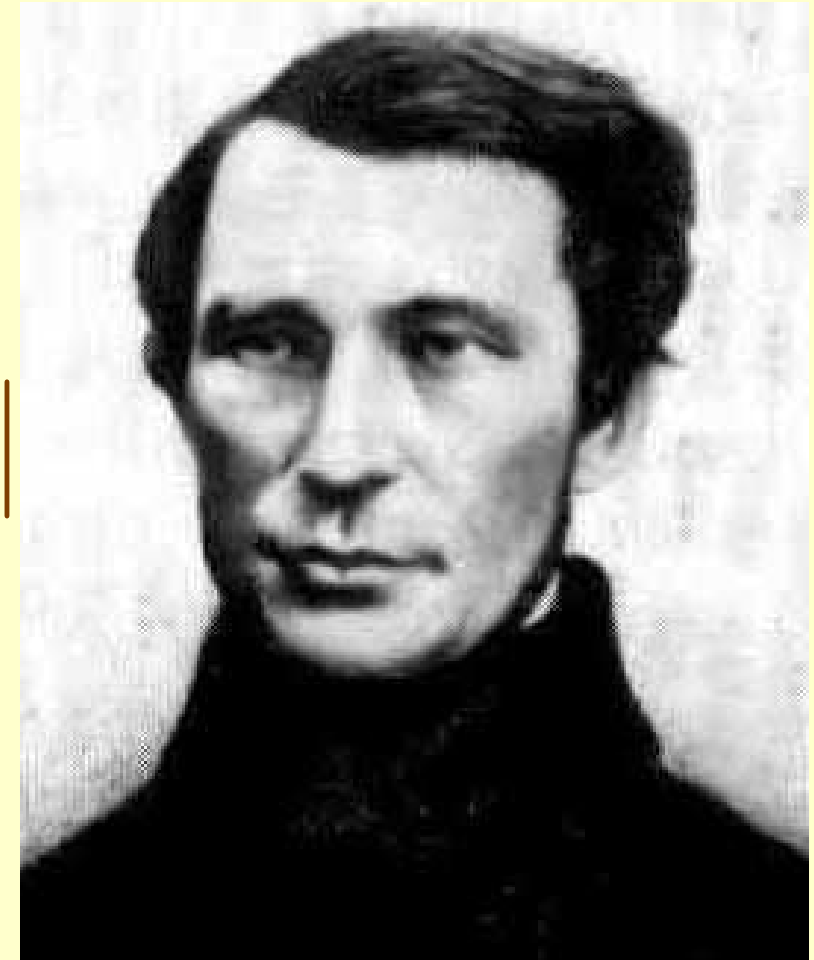
$$p_n(x) \equiv \sum_{k=0}^n f\left(\frac{k}{n}\right) \binom{n}{k} x^k (1-x)^{n-k}$$

Чебышев's theorem

- ▶ Чебышев: There is always a unique polynomial of any degree d which minimises

$$\|p - f\|_{\infty} = \max_{0 \leq x \leq 1} |p(x) - f(x)|$$

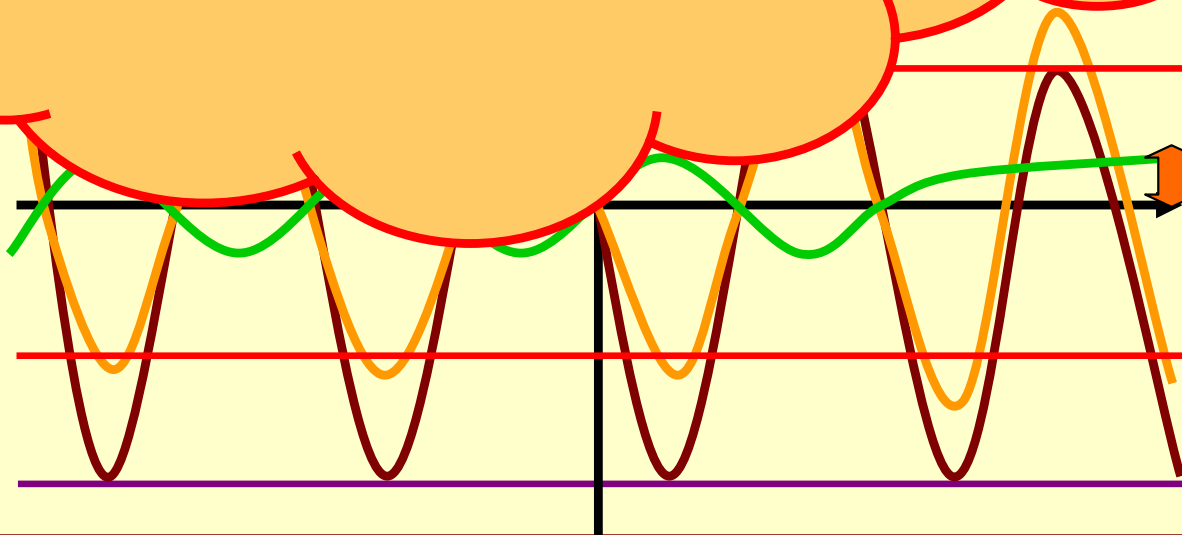
- ▶ The error $|p(x) - f(x)|$ reaches its maximum at at least $d+2$ points on the unit interval



Чебышев's theorem: Necessity

- ▶ Suppose $p-f$ has less than $d+2$ extrema of equal magnitude
- ▶ Then at most $d+1$ maxima exceed some magnitude
- ▶ This defines a
- ▶ Which has the opposite sign to
- ▶ (Lagrange's approximation)
- ▶ "to"
- ▶ A function than p to f

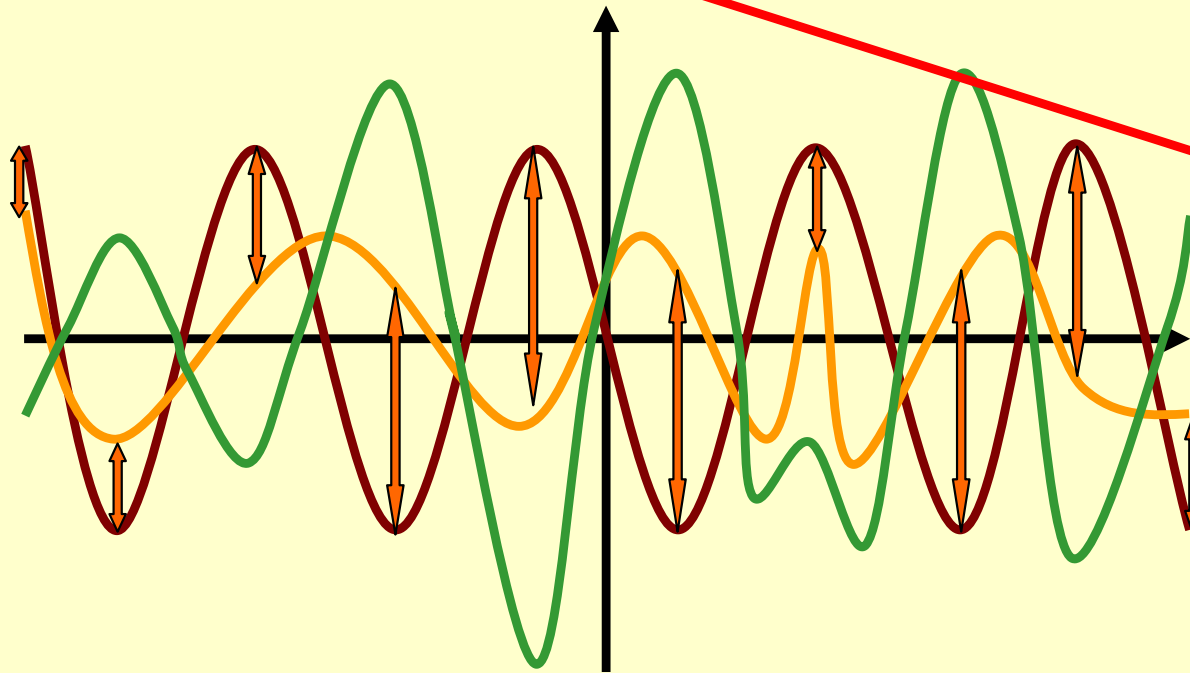
Lagrange was born in Turin!



Чебышев's theorem: Sufficiency

- ▶ Suppose there is a polynomial $\|p' - f\| < \|p - f\|$
- ▶ Then $|p'(x_i) - f(x_i)| \leq |p(x_i) - f(x_i)|$
- ▶ Therefore $p' - p$ must have $d+1$ zeros
- ▶ Thus $p' - p = 0$ as it is a polynomial of degree d

125th Anniversary of
Чебышев's birth





Чебышев polynomials

- ▶ Convergence is often exponential in d
 - ▶ The best approximation of degree $d-1$ over $[-1, 1]$ to x^d is $p_{d-1}(x) \equiv x^d - \left(\frac{1}{2}\right)^{d-1} T_d(x)$


- ▶ Where the Чебышев polynomials are

$$T_d(x) = \cos\left(d \cos^{-1}(x)\right)$$

- ▶ The notation is an old transliteration of Чебышев !
 - ▶ In general, truncated Чебышев L^2 approximation is not L^∞ optimal even for polynomials

- ▶ The error is $\|x^d - p_d(x)\|_\infty = \left(\frac{1}{2}\right)^{d-1} \|T_d(x)\|_\infty = 2e^{-d \ln 2}$

Чебышев rational functions

- ▶ Чебышев's theorem is easily extended to rational approximations
 - ▶ Rational functions with nearly equal degree numerator and denominator are usually best
 - ▶ Convergence is still often exponential
 - ▶ Rational functions usually give a much better approximation
- 
 - ▶ A simple (but somewhat slow) numerical algorithm for finding the optimal Чебышев rational approximation was given by Pemez



Чебышев rationals: Example

- ▶ A realistic example of a rational approximation is

$$\frac{1}{\sqrt{x}} \approx 0.3904603901 \frac{(x + 2.3475661045)(x + 0.1048344600)(x + 0.0073063814)}{(x + 0.4105999719)(x + 0.0286165446)(x + 0.0012779193)}$$

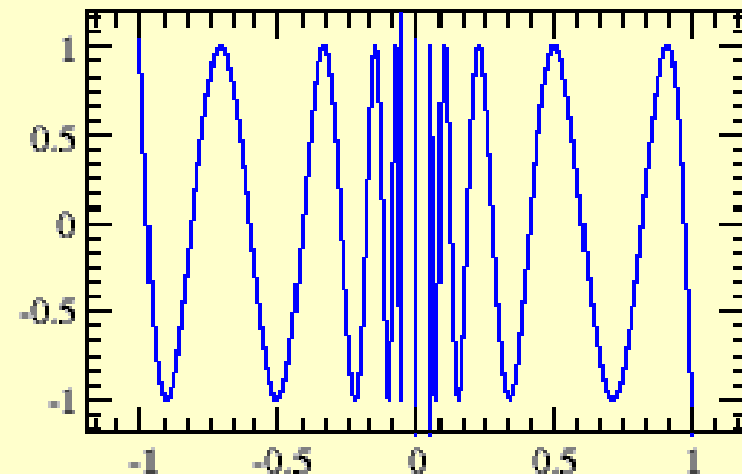
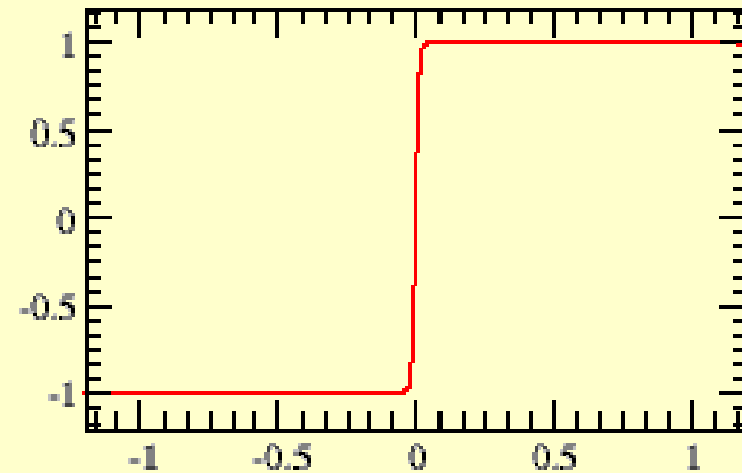
- ▶ This is accurate to within almost *0.1%* over the range $[0.003, 1]$
- ▶ Using a partial fraction expansion of such rational functions allows us to use a multishift linear equation solver, thus reducing the cost significantly.
- ▶ The partial fraction expansion of the rational function above is

$$\frac{1}{\sqrt{x}} \approx 0.3904603901 + \frac{0.0511093775}{x + 0.0012779193} + \frac{0.1408286237}{x + 0.0286165446} + \frac{0.5964845033}{x + 0.4105999719}$$

- ▶ This appears to be numerically stable.

Чебышев rationals: Example

- ▶ $\text{sgn}(x) = x R(x^2) + O(\Delta)$
- ▶ $\varepsilon < |x| < 1$
- ▶ $\varepsilon = 0.01$
- ▶ R of degree $(10, 10)$
- ▶ $\Delta = 0.00437985$
- ▶ Coefficients are known in closed form in terms of elliptic functions (Золотарев)



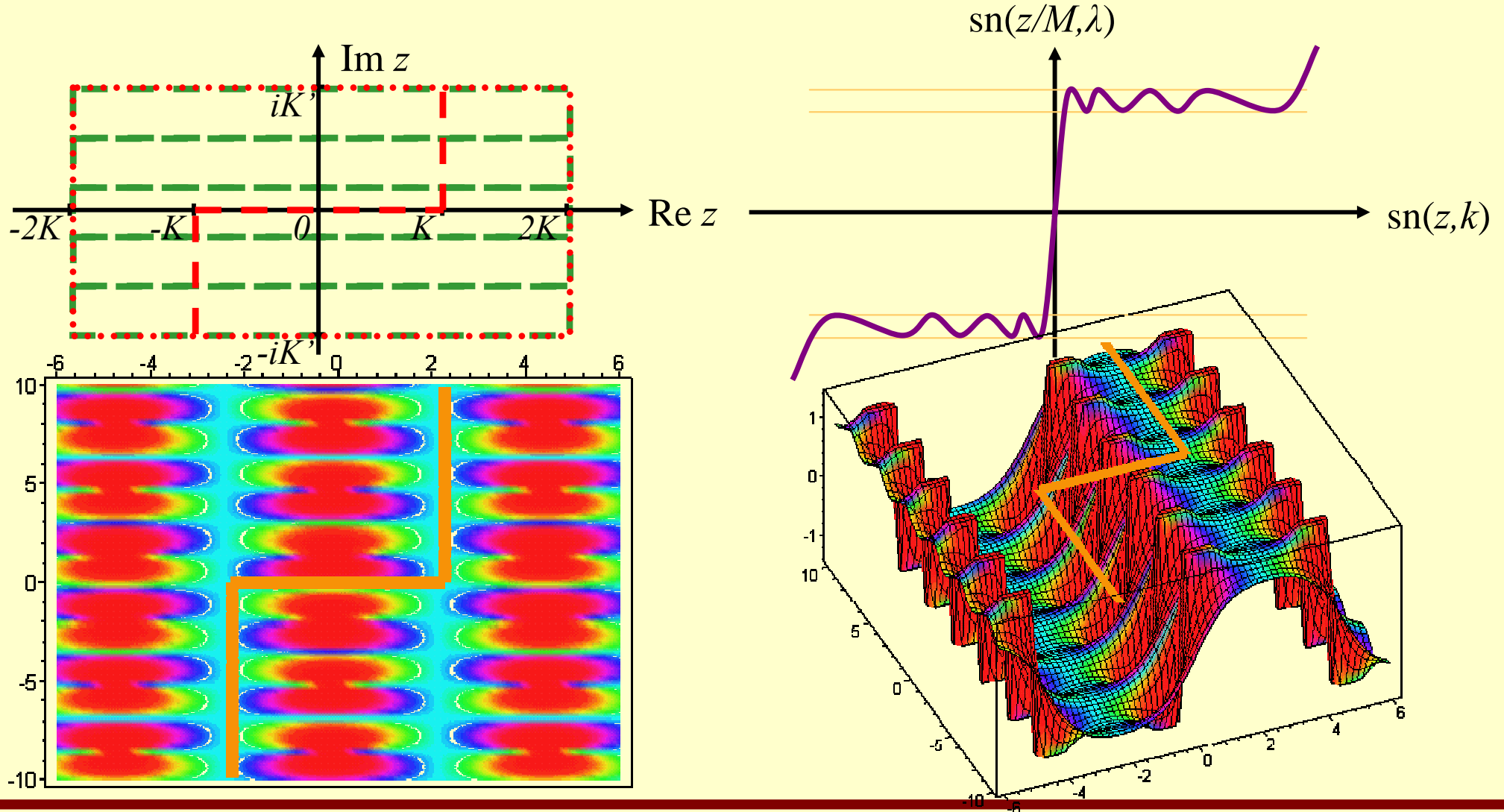
Золотарев's formula: I



- ▶ Modular transformation of Jacobi elliptic functions of degree n

$$\frac{\operatorname{sn}(z/M; \lambda)}{\operatorname{sn}(z; k)} = \frac{1}{M} \prod_{m=1}^{\lfloor n/2 \rfloor} \frac{1 - \frac{\operatorname{sn}(z; k)^2}{\operatorname{sn}(2iK'm/n; k)^2}}{1 - \frac{\operatorname{sn}(z; k)^2}{\operatorname{sn}(2iK'(m - \frac{1}{2})/n; k)^2}}$$

Золотарев's formula: II





Polynomials versus rationals

- ▶ Золотарев's formula has L_∞ error $\Delta \leq e^{\frac{n}{\ln \varepsilon}}$

- ▶ Optimal L^2 approximation with weight $\frac{1}{\sqrt{1-x^2}}$ is

$$\sum_{j=0}^n \frac{(-1)^j 4}{(2j+1)\pi} T_{2j+1}(x)$$

- ▶ This has L^2 error of $O(1/n)$
- ▶ Optimal L^∞ approximation cannot be too much better (or it would lead to a better L^2 approximation)

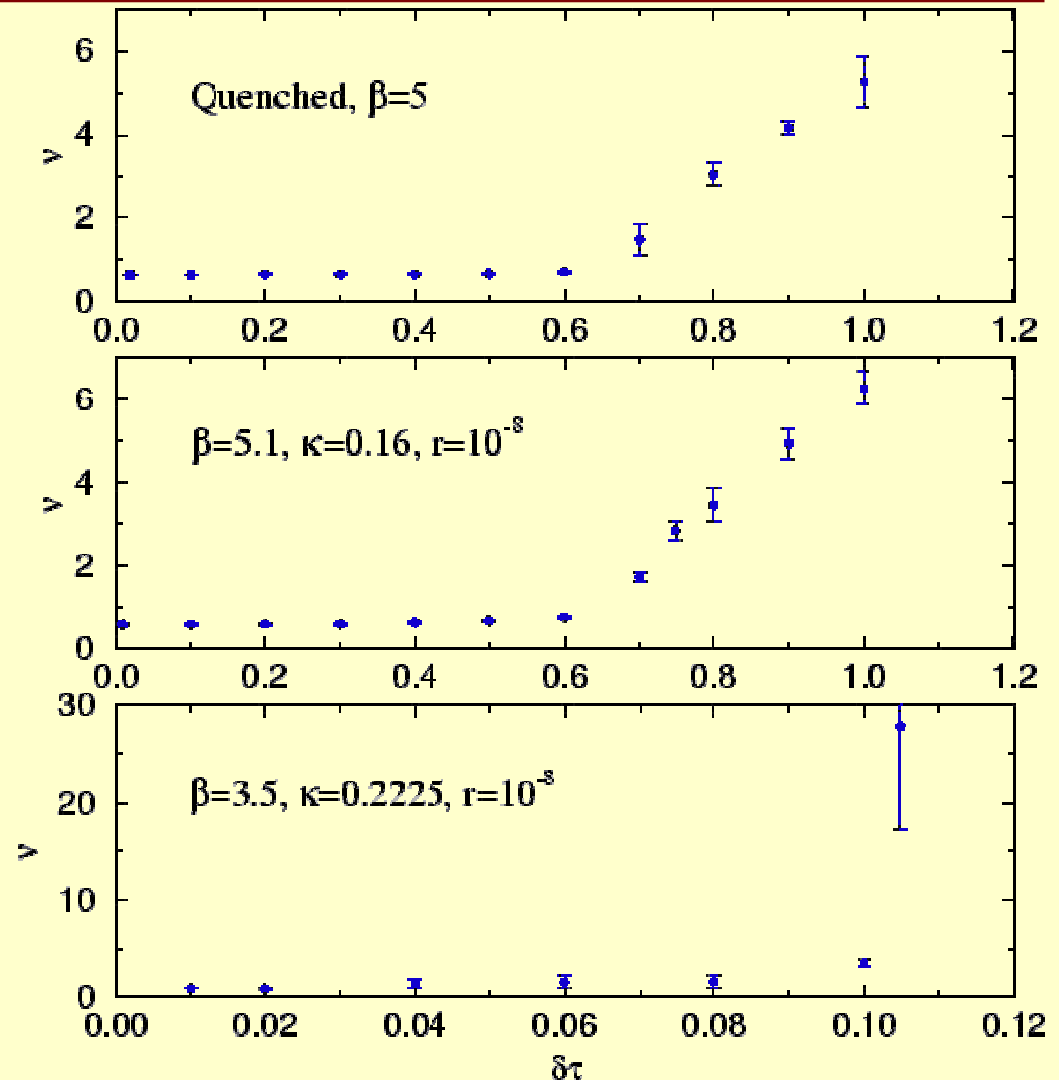


No Free Lunch Theorem

- ▶ We must apply the rational approximation with each CG iteration
 - ▶ $M^{1/n} \approx r(M)$
 - ▶ The condition number for each term in the partial fraction expansion is approximately $\kappa(M)$
 - ▶ So the cost of applying $M^{1/n}$ is proportional to $\kappa(M)$
 - ▶ Even though the condition number $\kappa(M^{1/n}) = \kappa(M)^{1/n}$
 - ▶ And even though $\kappa(r(M)) = \kappa(M)^{1/n}$
- ▶ So we don't win this way...

Instability of symplectic integrator

- ▶ Symplectic integrator
 - ▶ Leapfrog
 - ▶ Exactly reversible...
 - ▶ ...up to rounding errors
- ▶ Ляпунов exponent ν
 - ▶ $\nu > 0 \forall \delta\tau$
 - ▶ Chaotic equations of motion (Liu & Jansen)
 - ▶ $\nu \propto \delta\tau$ when $\delta\tau$ exceeds critical value $\delta\tau_c$
 - ▶ Instability of integrator
 - ▶ $\delta\tau_c$ depends on quark mass





Pseudofermions

- ▶ We want to evaluate a functional integral including the fermionic determinant $\det M$
- ▶ We write this as a bosonic functional integral over a pseudofermion field with kernel M^{-1}

$$\det M \propto \int d\phi^* d\phi e^{-\phi^* M^{-1} \phi}$$



Multipseudofermions

- ▶ We are introducing extra noise into the system by using a single pseudofermion field to sample this functional integral
 - ▶ This noise manifests itself as fluctuations in the force exerted by the pseudofermions on the gauge fields
 - ▶ This increases the maximum fermion force
 - ▶ This triggers the integrator instability
 - ▶ This requires decreasing the integration step size

- ▶ A better estimate is $\det M = [\det M^{1/n}]^n$

$$\det M^{\frac{1}{n}} \propto \int d\phi^* d\phi e^{-\phi^* M^{-\frac{1}{n}} \phi}$$



Hasenbusch's method

- ▶ Clever idea due to Hasenbusch
 - ▶ Start with the Wilson fermion action $M=1-\kappa H$
 - ▶ Introduce the quantity $M'=1-\kappa'H$
 - ▶ Use the identity $M = M'(M'^{-1}M)$
 - ▶ Write the fermion determinant as $\det M = \det M' \det (M'^{-1}M)$
 - ▶ Introduce separate pseudofermions for each determinant
 - ▶ Adjust κ' to minimise the cost
- ▶ Easily generalises
 - ▶ More than two pseudofermions
 - ▶ Wilson-clover action

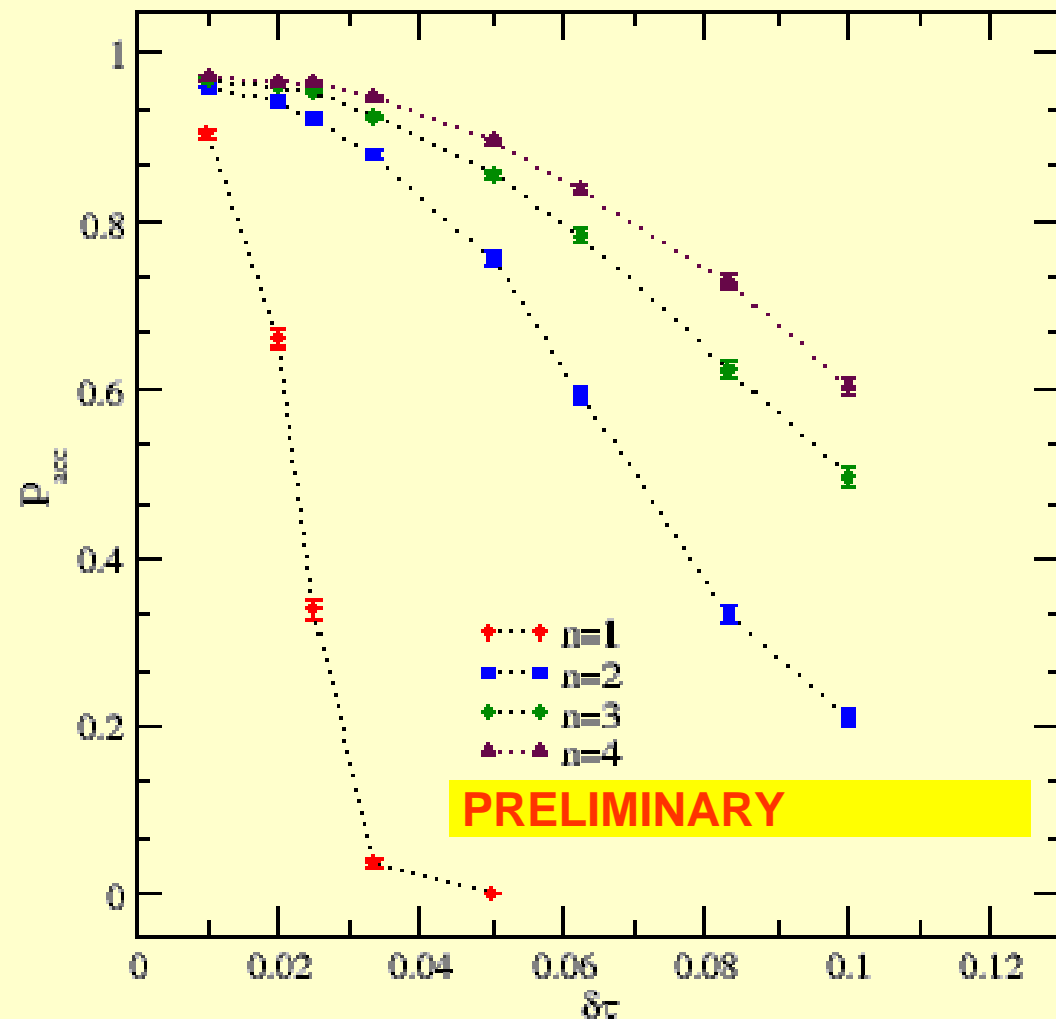


Violation of NFL Theorem

- ▶ So let's try using our n^{th} root trick to implement multipseudofermions
 - ▶ Condition number $\kappa(r(M)) = \kappa(M)^{1/n}$
 - ▶ So maximum force is reduced by a factor of $n\kappa(M)^{(1/n)-1}$
 - ▶ This is a good approximation if the condition number is dominated by a few isolated tiny eigenvalues
- ▶ Cost reduced by a factor of $n\kappa(M)^{(1/n)-1}$
 - ▶ Optimal value $n_{\text{opt}} \approx \ln \kappa(M)$
 - ▶ So optimal cost reduction is $(e \ln \kappa) / \kappa$
- ▶ This works!
- ▶ Numerical results due to Mike Clark

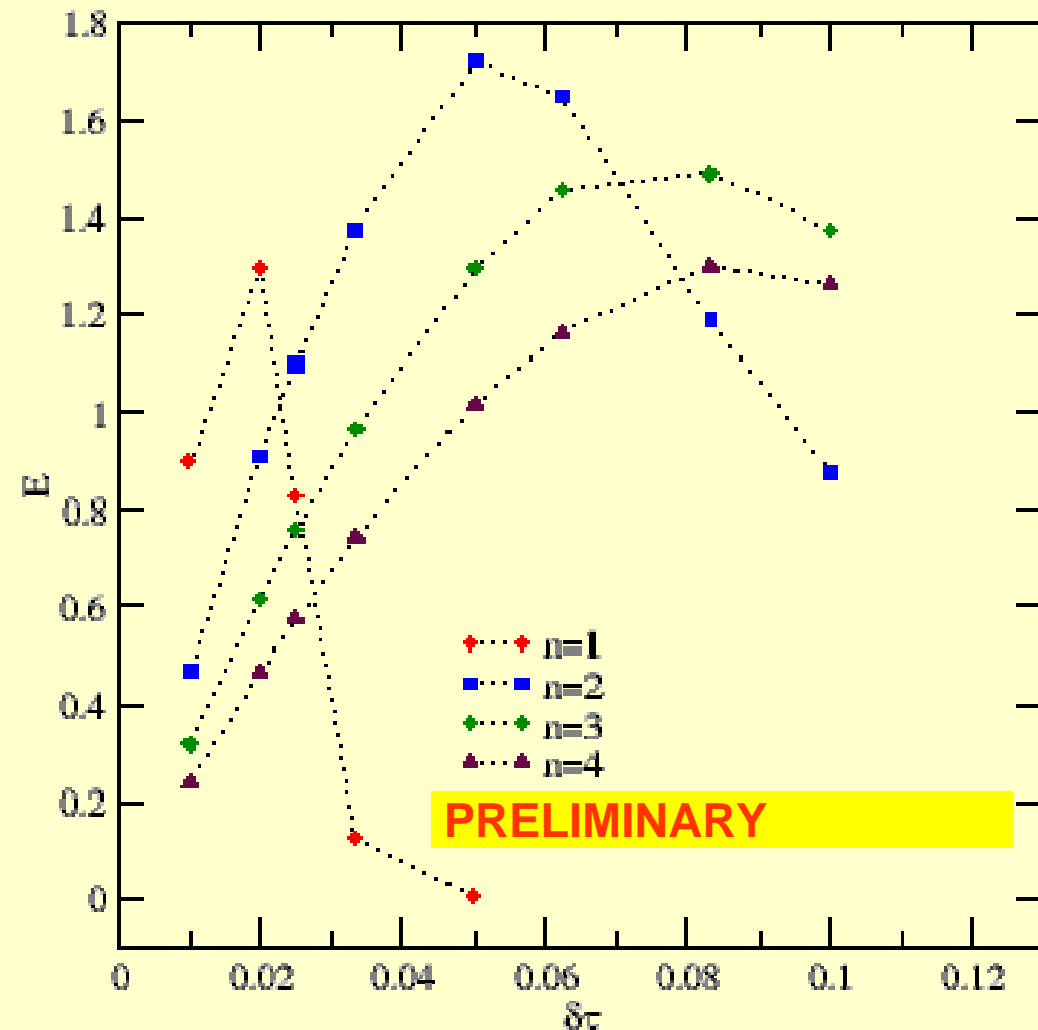
Acceptance rate at $m=0.025$

- ▶ 8^4 lattice
- ▶ $m=0.025$
- ▶ $\beta=5.26$
- ▶ $\tau=1.0$
- ▶ 4 flavours
- ▶ naïve staggered



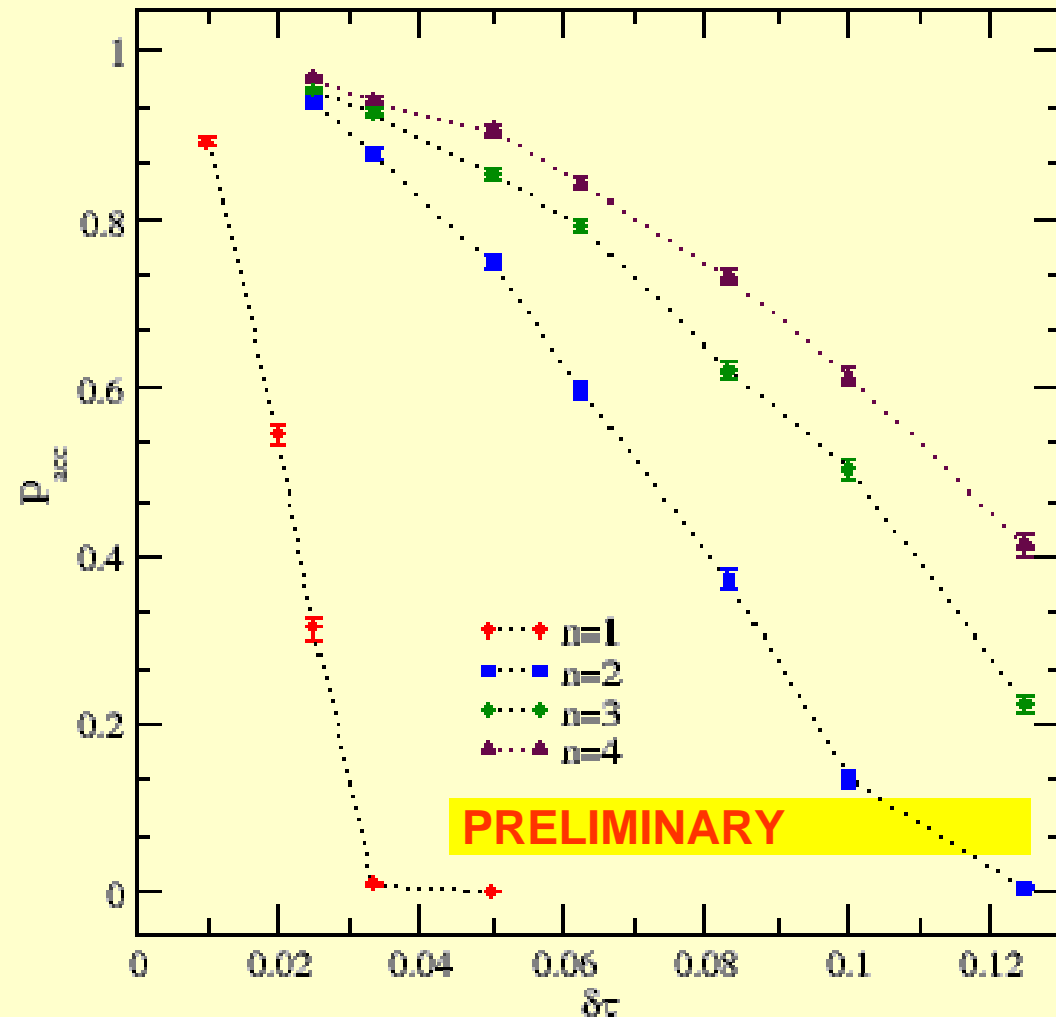
Efficiency at $m=0.025$

- ▶ 8^4 lattice
- ▶ $m=0.025$
- ▶ $\beta=5.26$
- ▶ $\tau=1.0$
- ▶ 4 flavours
- ▶ naïve staggered
- ▶ $n_{opt}=2$
- ▶ 33% gain



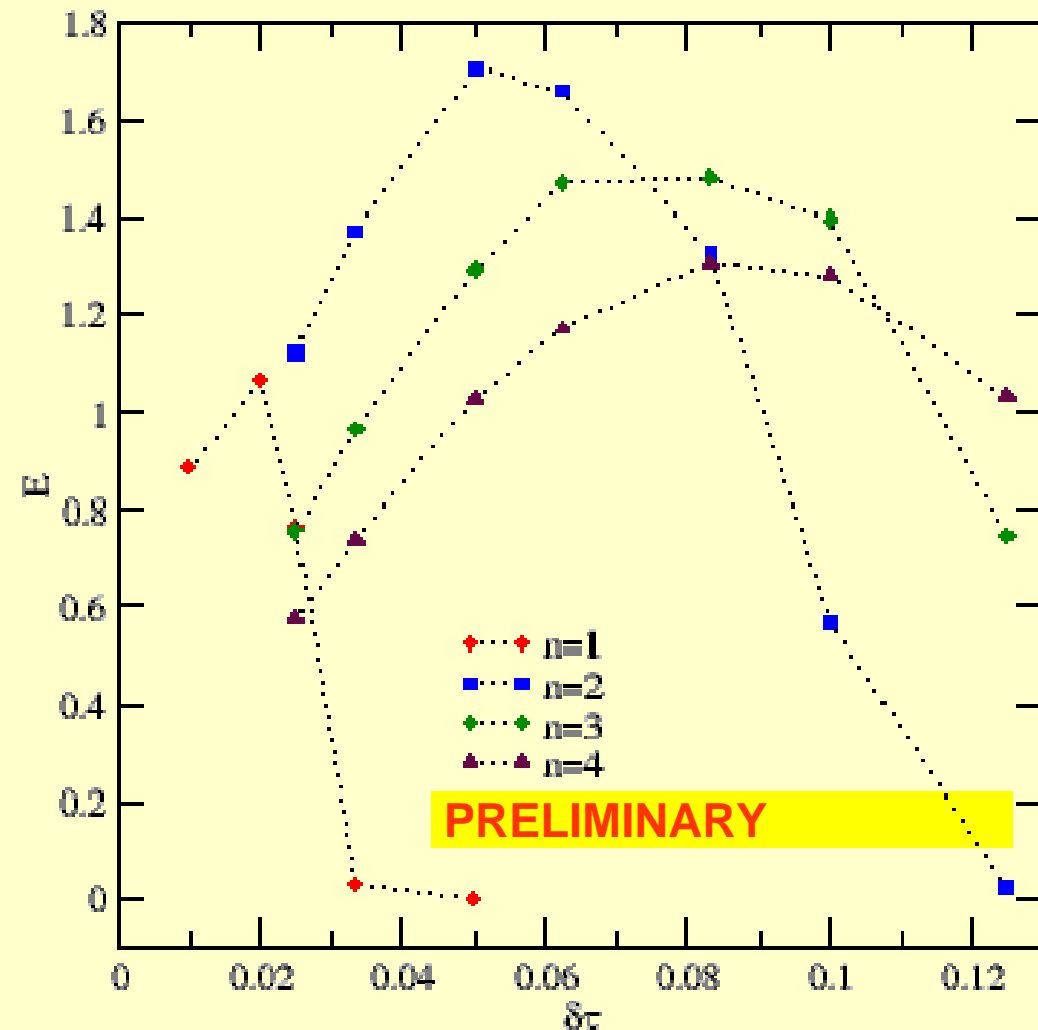
Acceptance rate at $m=0.01$

- ▶ 8^4 lattice
- ▶ $m=0.01$
- ▶ $\beta=5.26$
- ▶ $\tau=1.0$
- ▶ 4 flavours
- ▶ naïve staggered



Efficiency at $m=0.01$

- ▶ 8^4 lattice
- ▶ $m=0.01$
- ▶ $\beta=5.26$
- ▶ $\tau=1.0$
- ▶ 4 flavours
- ▶ naïve staggered
- ▶ $n_{opt}=2$
- ▶ 60% gain





RHMC technicalities

- ▶ In order to keep the algorithm exact to machine precision (as a Markov process)
 - ▶ Use a good (machine precision) rational approximation for the pseudofermion heatbath
 - ▶ Use a good rational approximation for the HMC acceptance test at the end of each trajectory
 - ▶ Use as poor a rational approximation as we can get away with to compute the MD force